

Set 9: Planning  
Classical Planning Systems  
Chapter 10 R&N

ICS 271 Fall 2017

# Outline: Planning

- Planning environments
- Classical Planning:
  - Situation calculus
  - PDDL: Planning domain definition language
- STRIPS Planning
- SAT planning
- Planning graphs
- Readings: Russel and Norvig chapter 10

# What is planning?

- “Planning is a task of finding a sequence of actions that will transfer the initial world into one in which the goal description is true.”
- “The planning can be seen as a sequence of actions generator which are restricted by constraints describing the limitations on the world under view.”
- “Planning as the process of devising, designing or formulating something to be done, such as the arrangements of the parts of a thing or an action or proceedings to be carried out.”

# Setup

- Actions : deterministic/non-deterministic?
- State variables : discrete/continuous?
- Current state : observable?
- Initial state : known?
- Actions : duration?
- Actions : 1 at a time?
- Objective : reach a goal? maximize utility/reward?
- Agent : 1 or more? Cooperative/competitive?
- Environment : Known/unknown, static?

# Setup

- Classical planning:
  - Actions : deterministic
  - States : fully observable, initial state known
  - Environment : known and static
  - Objective : reach a goal state
- Games
  - Agents : 2 (or more) competing
  - Objective : maximize utility
- Conformant planning:
  - Actions : non-deterministic
  - States : not observable, initial state unknown
  - Objective : maximize probability of reaching the goal
- Markov decision process (MDP):
  - Actions : non-deterministic with probabilities known
  - States : fully observable
  - Objective : maximize reward

# Planning vs Scheduling

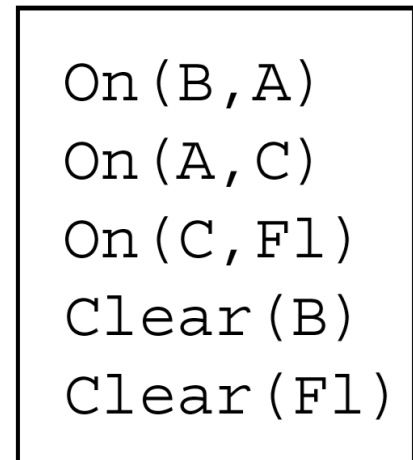
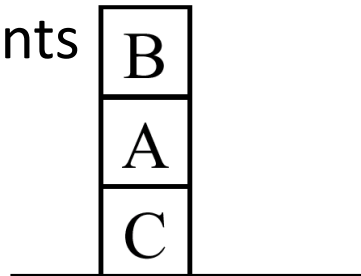
- Objective :
  - find a sequence of actions
  - find an allocation of jobs to resources
- Solution
  - Plan length unknown
  - Number of jobs to schedule known
- Complexity
  - PSPACE (planning)
  - NP-hard (scheduling)

# The Situation Calculus

- A **goal** can be described by a sentence:  
if we want to have a block on  $B$   $(\exists x)On(x, B)$
- **Planning**: finding a set of actions to achieve a goal sentence.
- **Situation Calculus** (McCarthy, Hayes, 1969, Green 1969)
  - A Predicate Calculus formalization of *states*, *actions*, and their *effects*.
  - $S_0$  state in the figure can be described by:

$On(B, A) \wedge On(A, C) \wedge On(C, Fl) \wedge Clear(B) \wedge clear(Fl)$

we reify the state and  
include them as arguments



# The Situation Calculus (continued)

- The atoms denotes relations over states called **fluents**.

$$On(B, A, S_0) \wedge On(A, C, S_0) \wedge On(C, Fl, S_0) \wedge clear(B, S_0)$$

- We can also have.

$$(\forall x, y, s)[On(x, y, s) \wedge \neg(y = Fl) \rightarrow \neg Clear(y, s)]$$

$$(\forall s)Clear(Fl, s)$$

- Knowledge about state and actions = predicate calculus knowledge base.
- Inference can be used to answer:
  - Is there a state satisfying a goal?
  - How can the present state be transformed into that state by actions? The answer is a *plan*



# Representing Actions

- Reify the actions: denote an action by a symbol
- actions are **functions**
- $\text{move}(B,A,\text{Floor})$ : move block A from block B to Floor
- $\text{move}(x,y,z)$  - action schema
- **do**: A function constant, **do** denotes a function that maps actions and states into states

$$\text{— } do(\alpha, \sigma) \rightarrow \sigma_1$$

action

state

# Representing Actions (continued)

- Express the effects of actions.
  - Example: (on, move) (expresses the effect of move on “On”)
  - Positive effect axiom:

$[On(x, y, s) \wedge Clear(x, s) \wedge Clear(z, s) \wedge (x \neq z) \rightarrow On(x, z, do(move(x, y, z), s))]$

- Negative effect axiom:

$[On(x, y, s) \wedge Clear(x, s) \wedge Clear(z, s) \wedge (x \neq z) \rightarrow \neg On(x, y, do(move(x, y, z), s))]$

- Positive: describes how action makes a fluent true
- Negative : describes how action makes a fluent false
- Antecedent: pre-condition for actions
- Consequent: how the fluent is changed

# Frame Axioms

- Not everything true can be inferred  
On(C,Fl) remains true but cannot be inferred
- Actions have local effect
  - We need **frame axioms** for each action and each fluent that does not change as a result of the action
  - example: frame axioms for (move, on)
  - If a block is on another block and *move* is not relevant, it will stay the same.

- Positive:

$$[On(x, y, s) \wedge (x \neq u)] \rightarrow On(x, y, do(move(u, v, z), s))$$

- Negative:

$$(\neg On(x, y, s) \wedge [(x \neq u) \vee (y \neq z)]) \rightarrow \neg On(x, y, do(move(u, v, z), s))$$

# Frame Axioms (continued)

- Frame axioms for (move, clear):

$$Clear(u, s) \wedge (u \neq z) \rightarrow Clear(u, do(move(x, y, z), s))$$

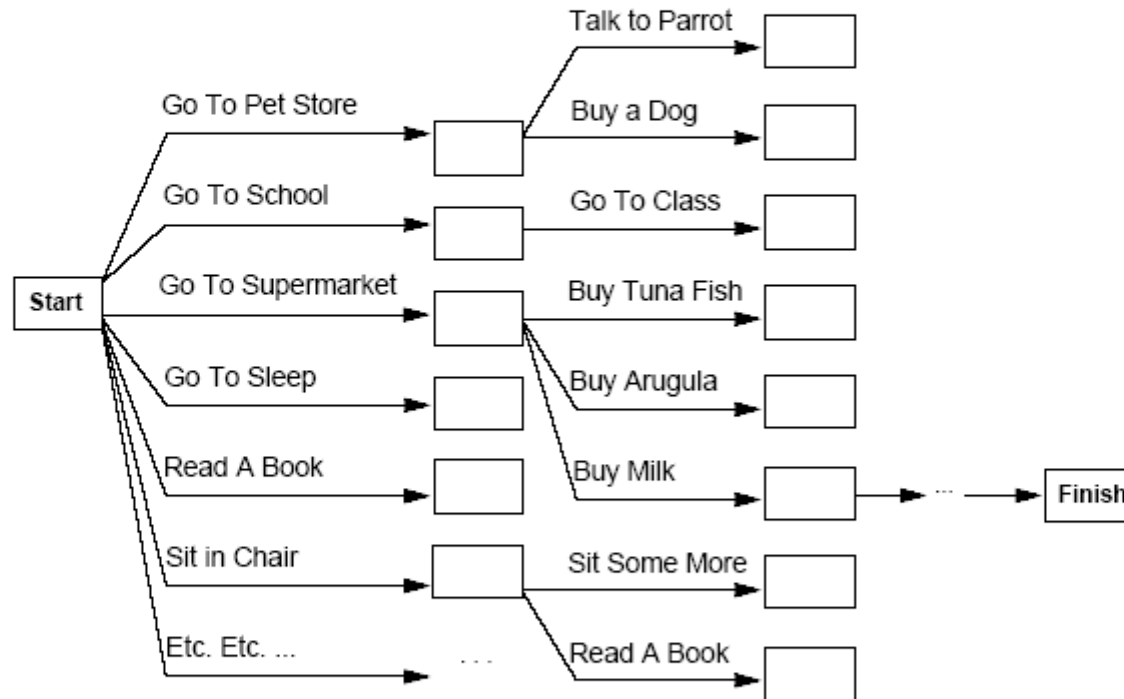
$$\neg Clear(u, s) \wedge (u \neq y) \rightarrow \neg Clear(u, do(move(x, y, z), s))$$

- **The frame problem:** need axioms for every combination of {action, predicate, fluent}!!!
- There are languages that embed some assumption on frame axioms that can be derived automatically:
  - Default logic
  - Negation as failure
  - Nonmonotonic reasoning
  - Minimizing change

## Search vs. planning

Consider the task *get milk, bananas, and a cordless drill*

Standard search algorithms seem to fail miserably:

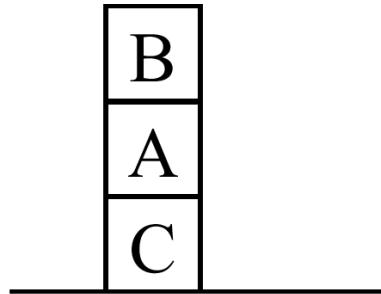


After-the-fact heuristic/goal test inadequate

PDDL: Planning Domain Definition  
Language  
STRIPS Planning systems

# STRIPS: describing goals and state

- On(B,A)
- On(A,C)
- On(C,Fl)
- Clear(B)
- Clear(Fl)



On (B , A)  
On (A , C)  
On (C , Fl )  
Clear (B)  
Clear (Fl)

- **State descriptions:** conjunctions of ground functionless atoms
  - **Factored representation of states!**
- A formula describes a set of world states :  $\text{On}(A,B) \wedge \text{Clear}(A)$
- Lifted version (schema):  $\text{On}(x,B) \wedge \text{Clear}(x)$
- Initial state is a conjunction of ground atoms
- Planning search for a formula satisfying a goal description
  - Goal wff:  $\exists x[P(x) \wedge Q(y)]$
  - Given a goal wff, the search algorithm looks for a sequence of **actions** that transforms initial state into a state description that entails the goal wff.

# STRIPS: description of actions

- A STRIPS operator (action) has 3 parts:
  - A set PC, of ground literals (*preconditions*)
  - A set D, of ground literals called the *delete list*
  - A set A, of ground literals called the *add list*
- Usually described by Schema: *Move(x,y,z)*
  - PC: *On(x,y)* and *Clear(x)* and *Clear(z)*
  - D: *Clear(z)*, *On(x,y)*
  - A: *On(x,z)*, *Clear(y)*, *Clear(Fl)*
- Lifting from prop logic level of representation to FOL level of representation
- A state  $S_{i+1}$  is created applying operator O by adding A and deleting D to/from  $S_i$ .



## STRIPS operators

Tidily arranged actions descriptions, restricted language

ACTION:  $Buy(x)$

PRECONDITION:  $At(p), Sells(p, x)$

EFFECT:  $Have(x)$

[Note: this abstracts away many important details!]

Restricted language  $\Rightarrow$  efficient algorithm

Precondition: conjunction of positive literals

Effect: conjunction of literals

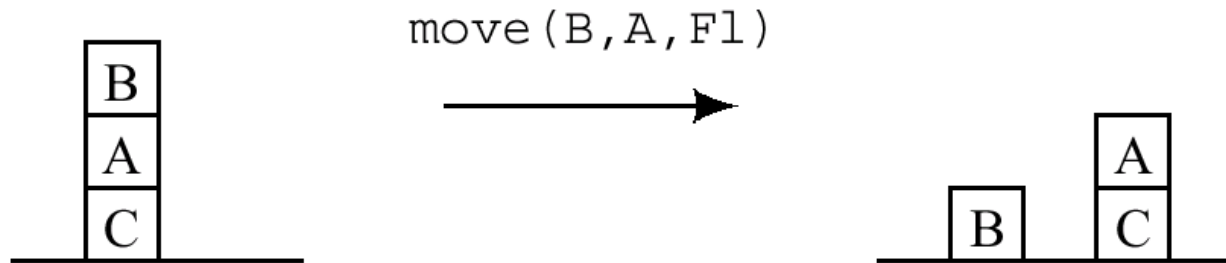
A complete set of STRIPS operators can be translated into a set of successor-state axioms

$At(p) Sells(p, x)$

**Buy(x)**

$Have(x)$

# Example: the move operator

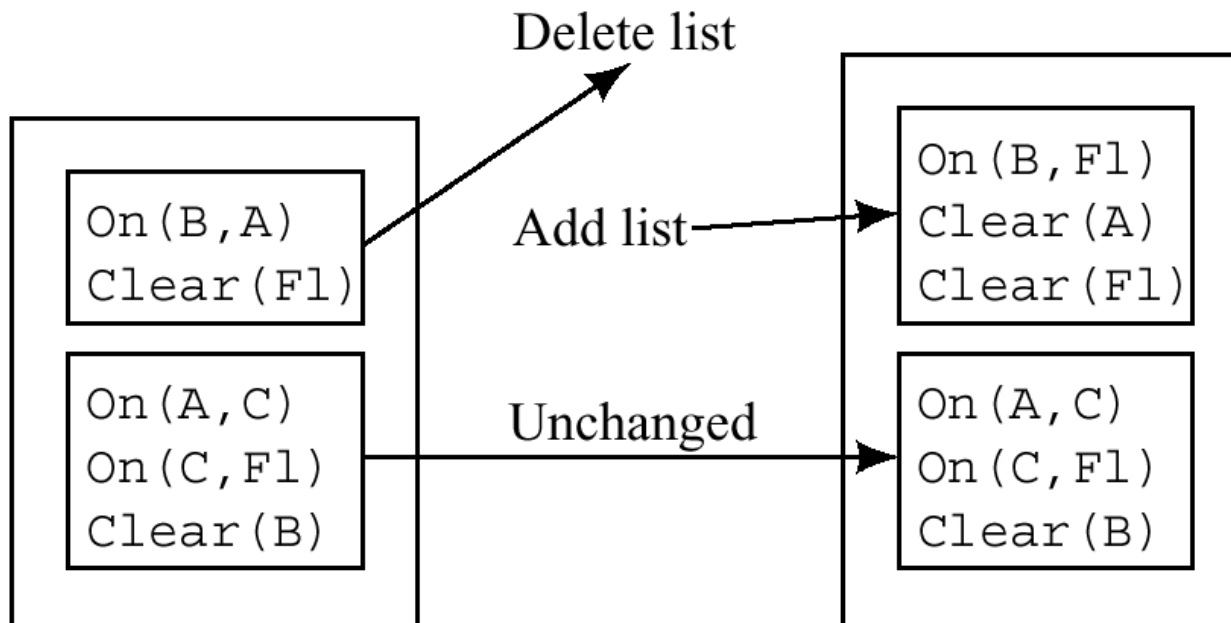


Precondition:

On (B, A)

Clear (B)

Clear (F1)



# PDDL vs STRIPS

- A language that yields a search problem : actions translate into operators in search space
- *PDDL is a slight generalization of STRIP language*
- A state is
  - a set of positive ground literals (STRIPS)
  - a set of ground literals (PDDL)
- Closed world assumption : fluents that are not mentioned are false (STRIPS).
- If a literals is not mentioned, it is unknown (PDDL).

- Action schema:

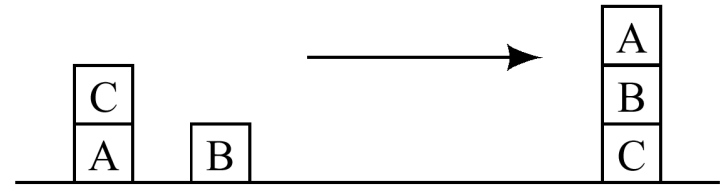
*Action(Fly(p,from,to)):*

*Precond:  $At(p,from) \wedge Plane(p) \wedge Airport(from) \wedge Airport(to)$*

*Effect:  $\neg At(p,from) \wedge At(p,to)$*

- *The schema consists of precondition and effect lists :*
  - *Only positive preconditions (STRIPS)*
  - *Positive or negative preconditions (PDDL)*
- *A set of action schemas is a definition of a planning domain.*
- *A specific problem is defined by an initial state (a set of ground atoms) and a goal: conjunction of atoms, some not grounded ( $At(p,SFO), Plane(p)$ )*

# The block world



```
Init(On(A, Table) ∧ On(B, Table) ∧ On(C, Table)
    ∧ Block(A) ∧ Block(B) ∧ Block(C)
    ∧ Clear(A) ∧ Clear(B) ∧ Clear(C))
Goal(On(A, B) ∧ On(B, C))
Action(Move(b, x, y),
    PRECOND: On(b, x) ∧ Clear(b) ∧ Clear(y) ∧ Block(b) ∧
        {b ≠ x} ∧ {b ≠ y} ∧ {x ≠ y},
    EFFECT: On(b, y) ∧ Clear(x) ∧ ¬ On(b, x) ∧ ¬ Clear(y))
Action(MoveToTable(b, x),
    PRECOND: On(b, x) ∧ Clear(b) ∧ Block(b) ∧ {b ≠ x},
    EFFECT: On(b, Table) ∧ Clear(x) ∧ ¬ On(b, x))
```

**Figure 11.4** A planning problem in the blocks world: building a three-block tower. One solution is the sequence  $[Move(B, Table, C), Move(A, Table, B)]$ .

# Summary so far

- Planning as inference : situation calculus
  - States defined by FOL sentences
  - Action effect sentences as FOL sentences
  - Frame axioms : for every action X predicate X object, define what effect non-related action has, as FOL sentences
  - Computational issues : ineffective inf procedure, semi-decidability of FOL
- Planning as search
  - PDDL (STRIPS) language
  - States defined by a set of literals (pos or neg)
  - Actions defined by action schemas : PC, AL/DL (Effects list)
  - An action can be executed in a state if PC is satisfied in the state
  - A set of action schemas = planning domain
  - Planning domain + initial/goal states = planning problem instance
  - This formulation naturally defines a search space
  - This formulation also lends itself to automatic heuristic generation

# A STRIP/PDDL description of an aircargo transportation problem

Problem: flying cargo in planes from one location to another

```
Init{At{C1, SFO} ∧ At{C2, JFK} ∧ At{P1, SFO} ∧ At{P2, JFK}  
  ∧ Cargo{C1} ∧ Cargo{C2} ∧ Plane{P1} ∧ Plane{P2}  
  ∧ Airport{JFK} ∧ Airport{SFO}}  
Goal{At{C1, JFK} ∧ At{C2, SFO}}  
Action{Load{c, p, a},  
  PRECOND: At{c, a} ∧ At{p, a} ∧ Cargo{c} ∧ Plane{p} ∧ Airport{a}  
  EFFECT: ¬ At{c, a} ∧ In{c, p}}  
Action{Unload{c, p, a},  
  PRECOND: In{c, p} ∧ At{p, a} ∧ Cargo{c} ∧ Plane{p} ∧ Airport{a}  
  EFFECT: At{c, a} ∧ ¬ In{c, p}}  
Action{Fly{p, from, to},  
  PRECOND: At{p, from} ∧ Plane{p} ∧ Airport{from} ∧ Airport{to}  
  EFFECT: ¬ At{p, from} ∧ At{p, to}}
```

**Figure 11.2** A STRIPS problem involving transportation of air cargo between airports.

*In*(*c*,*p*)- cargo *c* is inside plane *p*

*At*(*x*,*a*) – object *x* is at airport *a*

# STRIP for spare tire problem

Problem: Changing a flat tire

```
Init(At{Flat, Axle}  $\wedge$  At{Spare, Trunk})  
Goal(At{Spare, Axle})  
Action(Remove{Spare, Trunk},  
  PRECOND: At{Spare, Trunk}  
  EFFECT:  $\neg$  At{Spare, Trunk}  $\wedge$  At{Spare, Ground})  
Action(Remove{Flat, Axle},  
  PRECOND: At{Flat, Axle}  
  EFFECT:  $\neg$  At{Flat, Axle}  $\wedge$  At{Flat, Ground})  
Action(PutOn{Spare, Axle},  
  PRECOND: At{Spare, Ground}  $\wedge$   $\neg$  At{Flat, Axle}  
  EFFECT:  $\neg$  At{Spare, Ground}  $\wedge$  At{Spare, Axle})  
Action(LeaveOvernight,  
  PRECOND:  
  EFFECT:  $\neg$  At{Spare, Ground}  $\wedge$   $\neg$  At{Spare, Axle}  $\wedge$   $\neg$  At{Spare, Trunk}  
           $\wedge$   $\neg$  At{Flat, Ground}  $\wedge$   $\neg$  At{Flat, Axle})
```

**Figure 11.3** The simple spare tire problem.

# Complexity of classical planning

- Tasks
  - PlanSAT = decide if plan exists
  - Bounded PlanSAT = decide if plan of given length exists
- (Bounded) PlanSAT decidable but PSPACE-hard
- Disallow neg effects, (Bounded) PlanSAT NP-hard
- Disallow neg preconditions, PlanSAT in P but finding optimal (shortest) plan still NP-hard



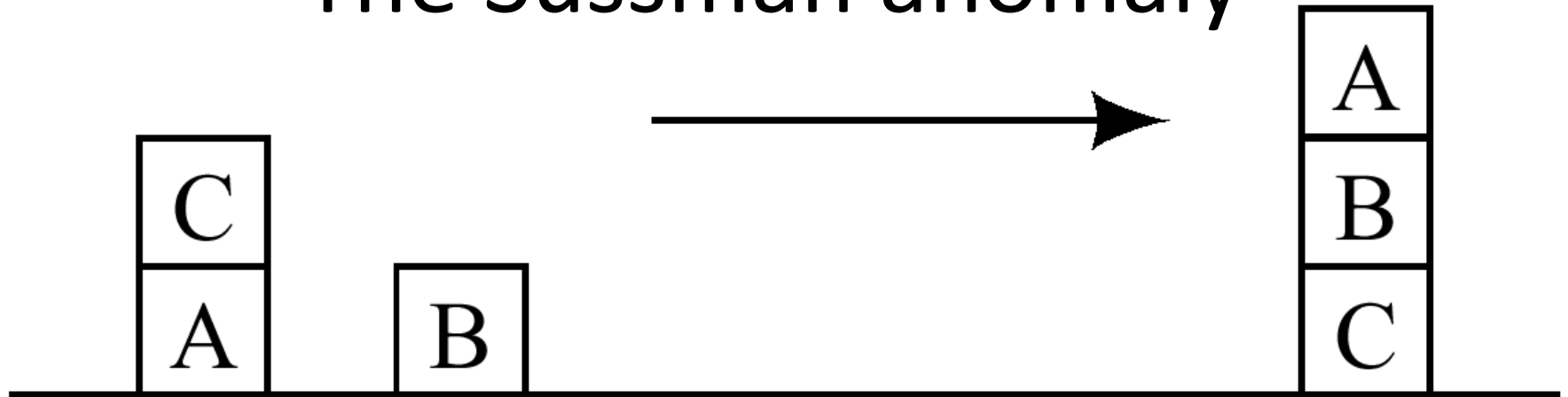
# Recursive STRIPS

- **STRIPS** algorithm :
  - Divide-and-Conquer forward search with islands
  - Achieve one subgoal at a time : achieve a new goal literal without ever violating already achieved goal literals or maybe temporarily violating previous subgoals.
- Motivated by **General Problem Solver (GPS)** by Newell Shaw and Simon (1959) - **Means-Ends** analysis.
- Each subgoal is achieved via a matched rule, then its preconditions are subgoals and so on. This leads to a planner called STRIPS( $\gamma$ ) when  $\gamma$  is a goal formula.

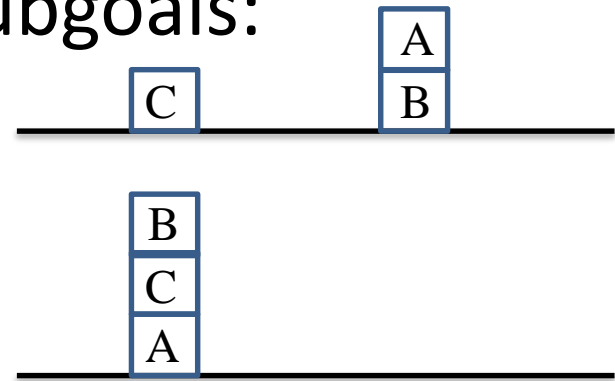
# Recursive STRIPS algorithm

- Algorithm maintains a set of goals
  - Start with all problem instance goals
  - At each iterations, take and satisfy one goal
- Algorithm :
  1. Take a goal from goal set
  2. Find a sequence of actions satisfying the goal from the current state, apply the actions, resulting in a new state.
  3. If stack empty, then done.
  4. Otherwise, the next goal is considered from the new state.
  5. At the end, check goals again.

# The Sussman anomaly



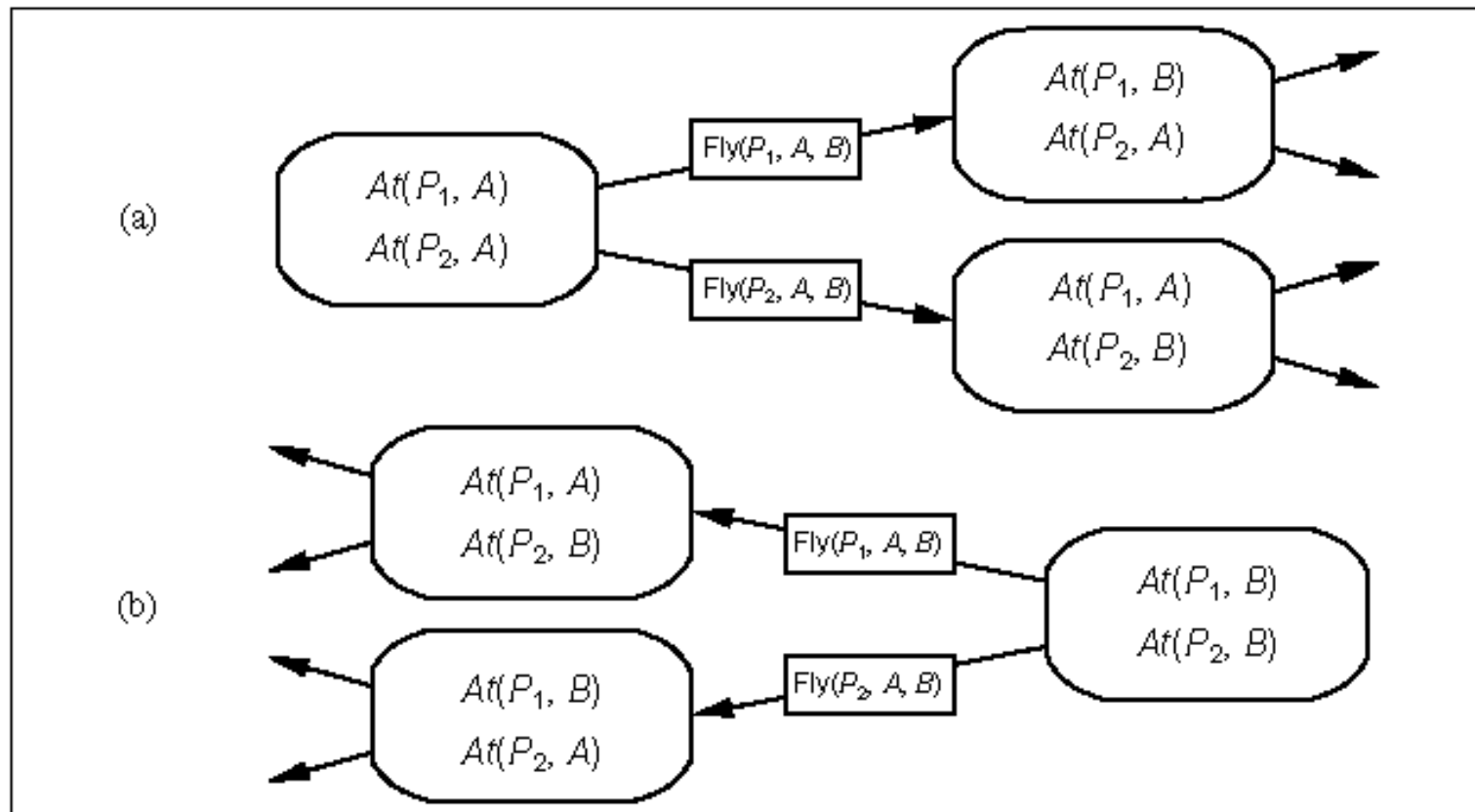
- RSTRIPS cannot find a valid plan
- Two possible orderings of subgoals:
  - $\text{On}(A,B)$  and  $\text{On}(B,C)$  or
  - $\text{On}(B,C)$  and  $\text{On}(A,B)$
- Non-interleaved planning does not work if goals are dependent



# Algorithms for Planning as State-space Search

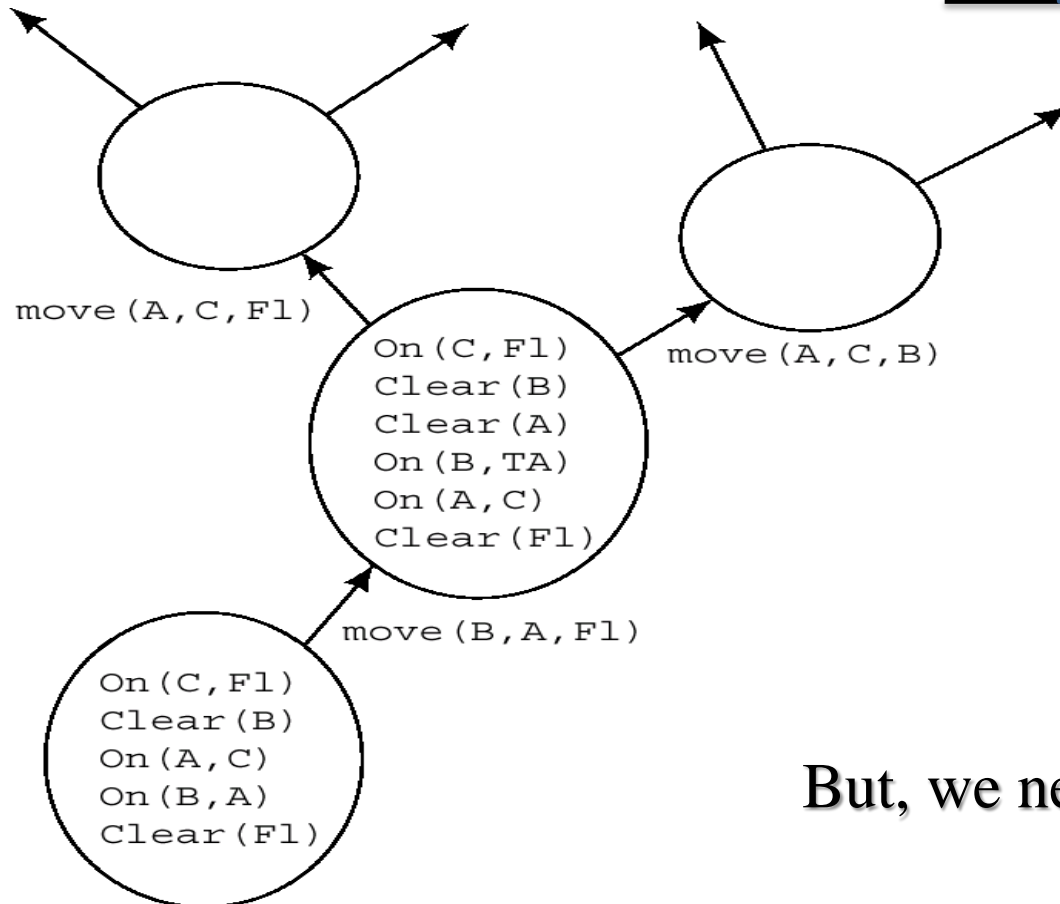
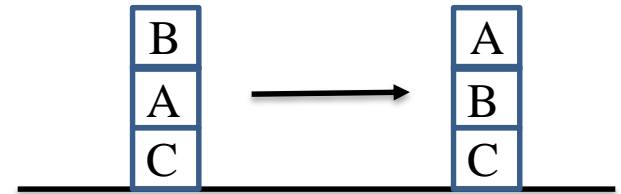
- Forward (progression) state-space search
  - Search with applicable actions
- Backward (regression) state-space search
  - Search with relevant actions
- Heuristic search
- Planning graphs
- Planning as satisfiability

# Planning forward and backward



**Figure 11.5** Two approaches to searching for a plan. (a) Forward (progression) state-space search, starting in the initial state and using the problem's actions to search forward for the goal state. (b) Backward (regression) state-space search: a belief-state search (see page 84) starting at the goal state(s) and using the inverse of the actions to search backward for the initial state.

# Forward Search Methods: can use A\* with some h and g

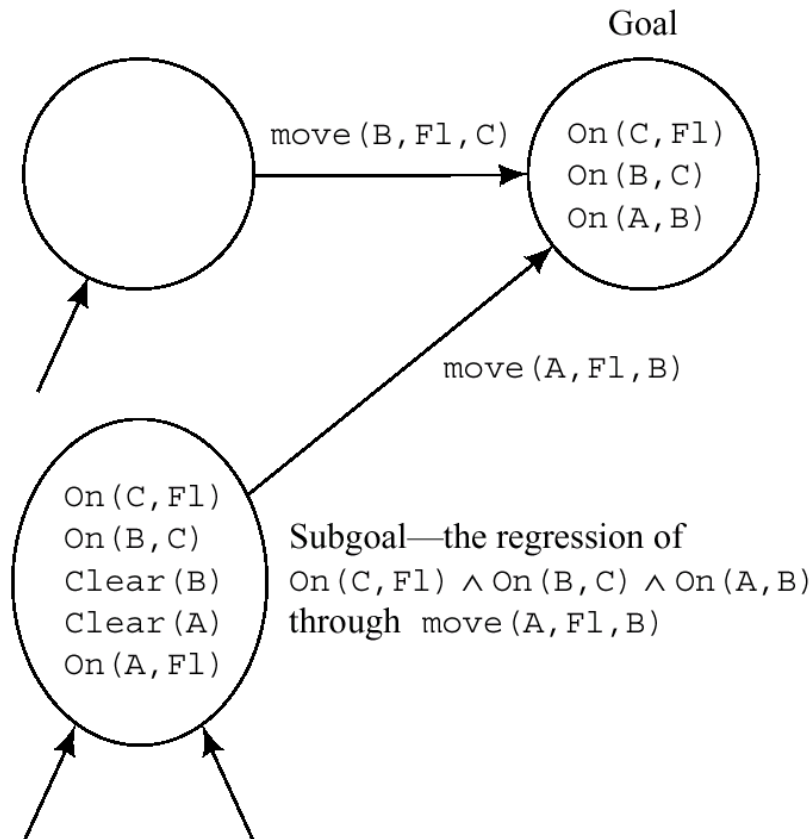


But, we need good heuristics

# Backward search methods

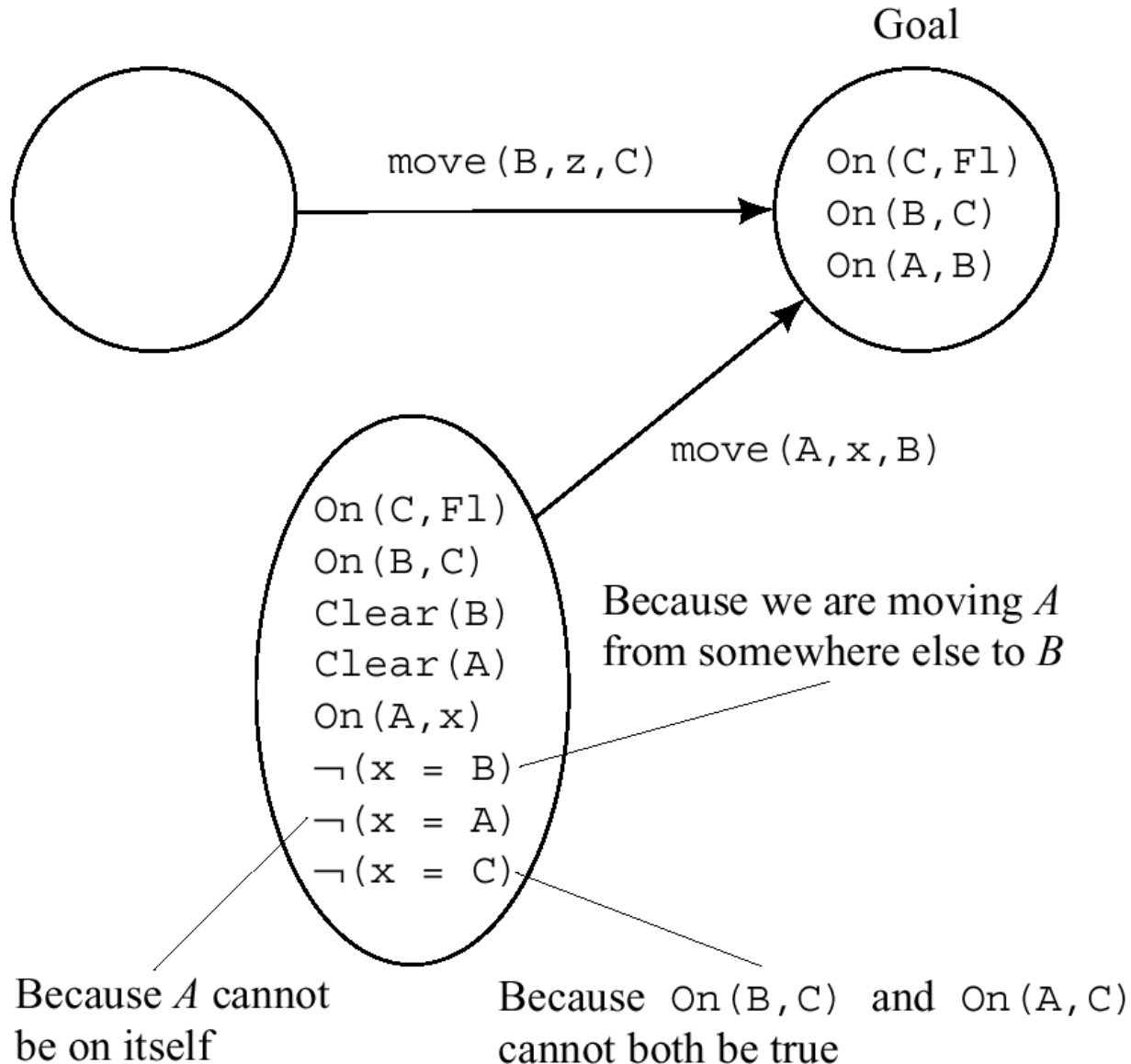
- Regressing a ground operator :

$$g' = (g - \text{ADD}(a)) \cup \text{PreCond}(a)$$



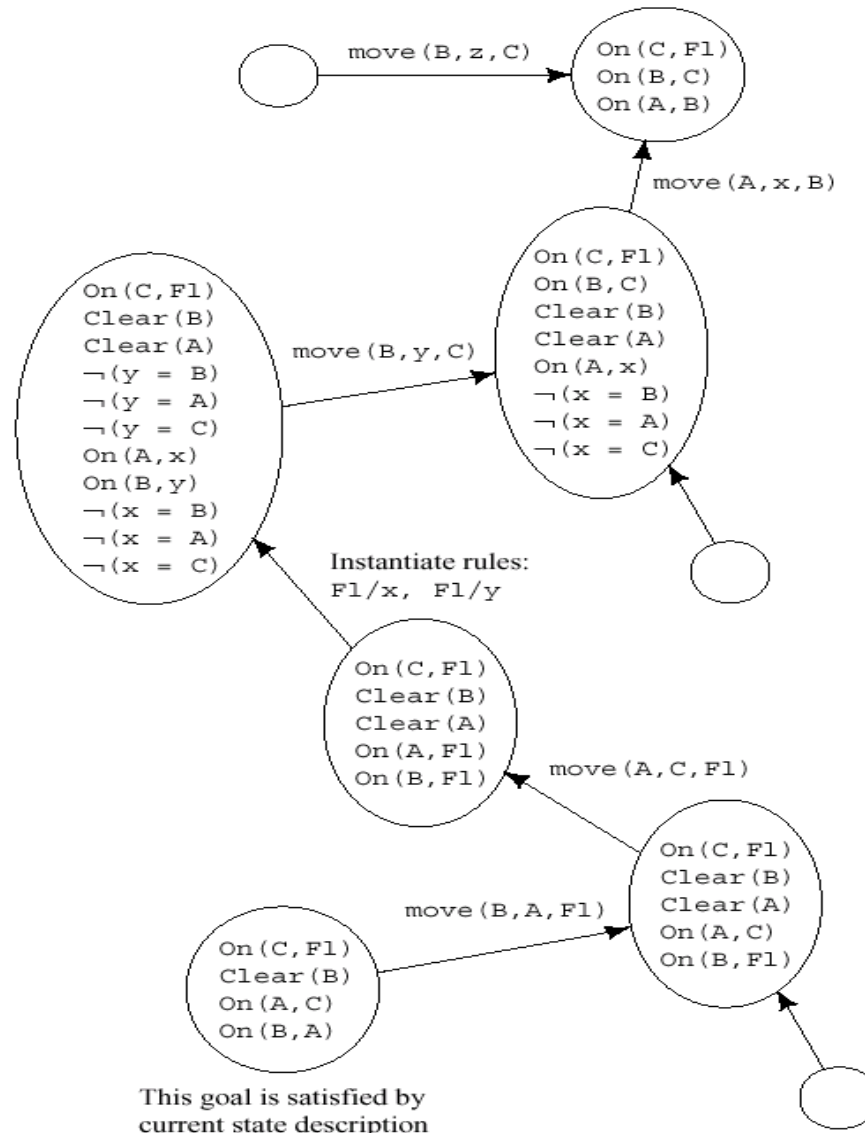
Continue until a subgoal is produced that is satisfied by current world state

# Regressing an action schema





# Example of Backward Search



# Forward vs Backward planning search

- Forward search space nodes correspond to individual (grounded) states of the plan state-space
- Backward search space nodes correspond to sets of plan state-space states, due to un-instantiated variables
  - because of this, designing good heuristics is hard(er)
  - however, it has smaller branching factor than FS
- Forward search only feasible if good heuristics available

# Heuristics for planning

- **Use relax problem idea** to get lower bounds on least number of actions to the goal
  - Add edges to the plan state-space graph
    - E.g. remove all or some preconditions
  - State abstraction (combining states)
- **Sub-goal independence**: compute the cost of solving each subgoal in isolation, and combine the costs, e.g. the sum of costs of solving each, or max cost
  - Can be pessimistic (interacting sub-plans)
  - Can be optimistic (negative effects)
- Various ideas related to removing negative/positive effects/preconditions.

# More on heuristic generation

- Ignore pre-conditions (example, 15 puzzle) : still hard, approximation easy but may not be admissible
- Ignore delete list: allow making monotone progress toward the goal.
  - Still NP-hard for optimal solution, but hill-climbing algorithms find an approximate solution in polynomial time that is admissible
- Abstraction: Combines many states into a single one: E.g. ignore some fluents, pattern databases
- FF : Fast-forward planner (Hoffman 2005), a forward state-space planner with
  - ignore-delete-list based heuristic
  - using planning graph to compute heuristic value
  - greedy search

# Planning Graphs

- A planning graph consists of a sequence of levels that correspond to time-steps in the plan
- Level 0 is the initial state.
- Each level contains a set of literals and a set of actions
- Literals are those that could be true at the time step.
- Actions are those that their preconditions could be satisfied at the time step.
- Works only for propositional planning.

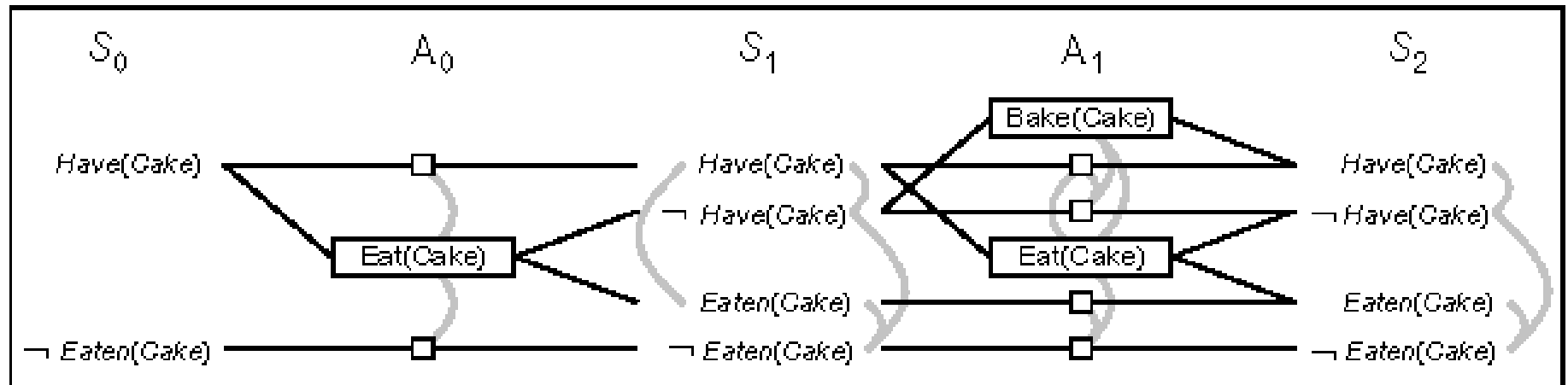
# Example: Have cake and eat it too

```
Init{Have{Cake}}  
Goal{Have{Cake}  $\wedge$  Eaten{Cake}}  
Action{Eat{Cake}}  
  PRECOND: Have{Cake}  
  EFFECT:  $\neg$  Have{Cake}  $\wedge$  Eaten{Cake}}  
Action{Bake{Cake}}  
  PRECOND:  $\neg$  Have{Cake}}  
  EFFECT: Have{Cake}}
```

**Figure 11.11** The “have cake and eat cake too” problem.

# The Planning graphs for “have cake”,

- Persistence actions: Represent “inactions” by boxes: frame axiom
- Mutual exclusions (mutex) are represented between literals and actions.
- $S_1$  represents multiple states
- Continue until two levels are identical. The graph levels off.
- The graph records the impossibility of certain choices using mutex links.
- Complexity of graph generation: polynomial in number of literals.



**Figure 11.12** The planning graph for the “have cake and eat cake too” problem up to level  $S_2$ . Rectangles indicate actions (small squares indicate persistence actions) and straight lines indicate preconditions and effects. Mutex links are shown as curved gray lines.

# Defining Mutex relations

- A **mutex** relation holds between **2 actions** on the same level iff any of the following holds:
  - **Inconsistency effect:** one action negates the effect of another. Example “Eat(Cake) and persistence of Have(cake)”
  - **Interference:** One of the effects of one action is the negation of the precondition of the other. Example “Eat(Cake) and persistence of Have(cake)”
  - **Competing needs:** one of the preconditions of one action is mutually exclusive with a precondition of another. Example: Bake(cake) and Eat(Cake).
- A **mutex** relation holds between **2 literals** at the same level iff
  - one is the negation of the other or if each possible pair of actions that can achieve the 2 literals is mutually exclusive



# Properties of planning graphs: termination

- Literals increase monotonically
  - Once a literal is in a level it will persist to the next level
- Actions increase monotonically
  - Since the precondition of an action was satisfied at a level and literals persist the action's precondition will be satisfied from now on
- Mutexes decrease monotonically:
  - If two actions are mutex at level  $S_i$ , they will be mutex at all previous levels at which they both appear
  - If two literals are not mutex, they will always be non-mutex later
- Because literals increase and mutex decrease it is guaranteed that we will have a level where  $S_i = S_{i-1}$  and  $A_i = A_{i-1}$  that is, the planning graph has stabilized

# Planning graphs for heuristic estimation

- Estimate the cost of achieving a goal by the level in the planning graph where it appears.
- To estimate the cost of a conjunction of goals use one of the following:
  - Max-level: take the maximum level of any goal (admissible)
  - Sum-cost: Take the sum of levels (inadmissible)
  - Set-level: find the level where they all appear without Mutex (admissible). Dominates max-level.
- Note, we don't have to build planning graph to completion to compute heuristic estimates
- Graph plans are an approximation of the problem. Representing more than pair-wise mutex is not cost-effective
  - E.g.  $\text{On}(A,B)$ ,  $\text{On}(B,C)$ ,  $\text{On}(C,A)$

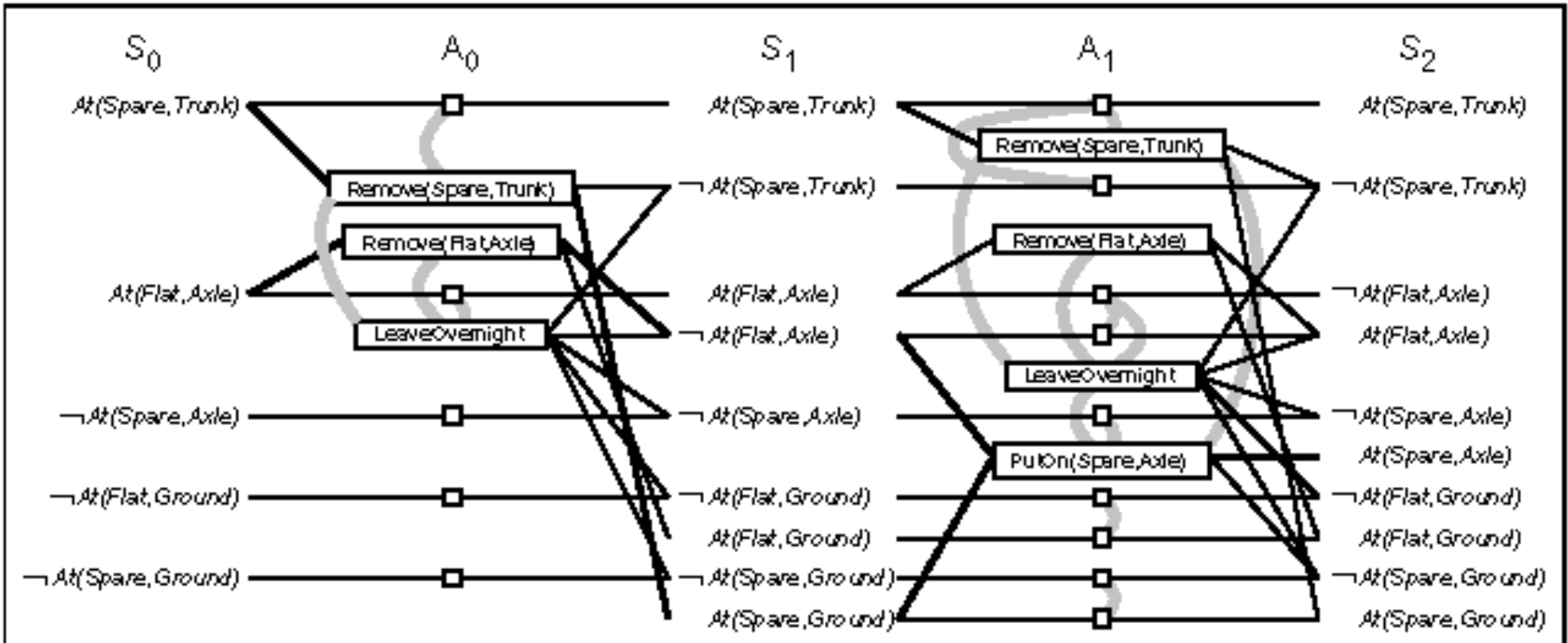
# The GraphPlan algorithm

- Start with a set of problem goals  $G$  at the last level  $S$
- At each level  $S_i$ , select a subset of conflict-free actions  $A_i$  for the goals of  $G_i$ , such that
  - Goals  $G_i$  are covered
  - No 2 actions in  $A_i$  are mutex
  - No 2 preconditions of any 2 actions in  $A_i$  are mutex
- Preconditions of  $A_i$  become goals of  $S_{i-1}$
- Success iff  $G_0$  is subset of initial state

# Planning graph for spare tire

goal:  $At(Spare, Axle)$

- $S_2$  has all goals and no mutex so we can try to extract solutions
- Use either CSP algorithm with actions as variables
- Or search backwards



**Figure 11.14** The planning graph for the spare tire problem after expansion to level  $S_2$ . Mutex links are shown as gray lines. Only some representative mutexes are shown, because the graph would be too cluttered if we showed them all. The solution is indicated by bold lines and outlines.

# The GraphPlan algorithm

```
function GRAPHPLAN(problem) returns solution or failure  
  
graph ← INITIAL-PLANNING-GRAPH(problem)  
goals ← GOALS[ problem ]  
loop do  
  if goals all non-mutex in last level of graph then do  
    solution ← EXTRACT-SOLUTION(graph, goals, LENGTH(graph))  
    if solution ≠ failure then return solution  
    else if NO-SOLUTION-POSSIBLE(graph) then return failure  
  graph ← EXPAND-GRAPH(graph, problem)
```

**Figure 11.13** The GRAPHPLAN algorithm. GRAPHPLAN alternates between a solution extraction step and a graph expansion step. EXTRACT-SOLUTION looks for whether a plan can be found, starting at the end and searching backwards. EXPAND-GRAPH adds the actions for the current level and the state literals for the next level.

# Searching planning-graph backwards with heuristics

- How to choose an action during backwards search:
  - Use greedy algorithm based on the level cost of the literals.
- For any set of goals:
  - 1. Pick first the literal with the highest level cost.
  - 2. To achieve the literal, choose the action with the easiest preconditions first (based on sum or max level of precondition literals).

# Main classical planning approaches

- The most effective approaches to planning currently are:
  - Forward state-space search with carefully crafted heuristics
  - Search using planning graphs (GraphPlan or CSP)
  - Translating to Boolean Satisfiability

# Planning as Satisfiability

- Index propositions with time steps:
  - $On(A,B)_0, ON(B,C)_0$
- Goal conditions:
  - the goal conjuncts at time  $t$ ,  $t$  is determined arbitrarily.
- Initial state :
  - Assert (pos) what is known, and (neg) what is not known.
- Actions: a proposition for each action for each time slot.
  - Exactly one action proposition is true are at  $t$  if serial plan is required
- Formula : if action is true, then effect must hold
- Formula : if action is true, then preconditions must have held
- Successor state axioms need to be expressed for each action (like in the situation calculus but it is propositional)
  - $F^{t+1} \Leftrightarrow ActionCausesF^t \vee (F^t \wedge \neg ActionCausesNotF^t)$



# Planning with propositional logic (continued)

- We write the formula:
  - initial state **and** action effect/precondition axioms **and** successor state axioms **and** goal state
- We search for a model to the formula. Those actions that are assigned true constitute a plan.
- To have a single plan we may have a mutual exclusion for all actions in the same time slot.
- We can also choose to allow partial order plans and only write exclusions between actions that interfere with each other.
- Planning: iteratively try to find longer and longer plans.

# SATplan algorithm

```
function SATPLAN(problem,  $T_{\text{max}}$ ) returns solution or failure
  inputs: problem, a planning problem
            $T_{\text{max}}$ , an upper limit for plan length

  for  $T = 0$  to  $T_{\text{max}}$  do
    cnf, mapping  $\leftarrow$  TRANSLATE-TO-SAT(problem,  $T$ )
    assignment  $\leftarrow$  SAT-SOLVER(cnf)
    if assignment is not null then
      return EXTRACT-SOLUTION(assignment, mapping)
  return failure
```

**Figure 11.15** The SATPLAN algorithm. The planning problem is translated into a CNF sentence in which the goal is asserted to hold at a fixed time step  $T$  and axioms are included for each time step up to  $T$ . (Details of the translation are given in the text.) If the satisfiability algorithm finds a model, then a plan is extracted by looking at those proposition symbols that refer to actions and are assigned *true* in the model. If no model exists, then the process is repeated with the goal moved one step later.

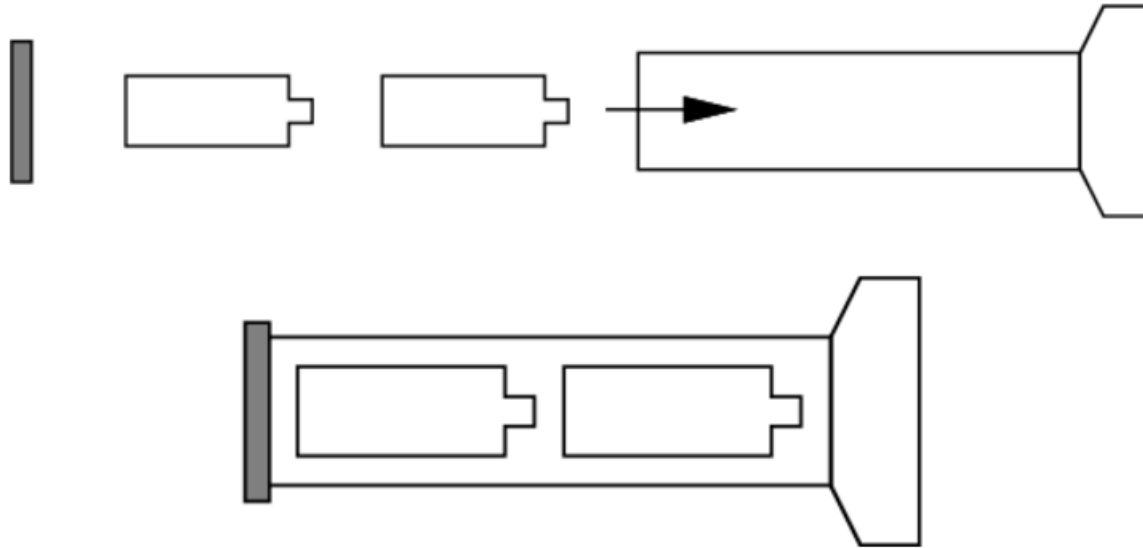
# Complexity of satplan

- The total number of action symbols is:
  - $|T| \times |Act| \times |O|^p$
  - $O$  = number of objects,  $p$  is scope of atoms.
- Number of clauses is higher.
- Example: 10 time steps, 12 planes, 30 airports, the complete action exclusion axiom has 583 million clauses.

# The flashlight problem (from Steve Lavelle)

- **Figure 2.18:** Three operators for the flashlight problem. Note that an operator can be expressed with variable argument(s) for which different instances (constants/objects) could be substituted.
- <http://planning.cs.uiuc.edu/node59.html#for:strips>
- Here is a SATplan for flashlight Battery
- <http://planning.cs.uiuc.edu/node68.html>

# Flashlight problem



- 4 objects : Cap, Battery1, Battery2, Flashlight
- 2 predicates : On (e.g.  $\text{On}(C,F)$ ), In (e.g.  $\text{In}(B1,F)$ )
- Initial state :  $\text{On}(C,F)$
- Assume initially : not  $\text{In}(B1,F)$  and not  $\text{In}(B2,F)$
- Goal :  $\text{On}(C,F)$ ,  $\text{In}(B1,F)$ ,  $\text{In}(B2,F)$

# Flashlight Problem

- 3 actions
  - PlaceCap
  - RemoveCap
  - Insert(i)

| Name             | Preconditions  | Effects                        |
|------------------|--|--------------------------------|
| <i>PlaceCap</i>  | $\{\neg On(Cap, Flashlight)\}$                         | $\{On(Cap, Flashlight)\}$      |
| <i>RemoveCap</i> | $\{On(Cap, Flashlight)\}$                              | $\{\neg On(Cap, Flashlight)\}$ |
| <i>Insert(i)</i> | $\{\neg On(Cap, Flashlight), \neg In(i, Flashlight)\}$ | $\{In(i, Flashlight)\}$        |

- Plan has 4 steps :
  - RemoveCap, Insert(B1), Insert(B2), PlaceCap

# SATPlan

- Guess length of plan **K**
- Initial state : conjunction of initial state literals and negation of all positive literals not given
- For each action and each time slot  $k$ 
  - $a_k \rightarrow (p_{k,1} \wedge \dots \wedge p_{k,m}) \wedge (e_{k+1,1} \wedge \dots \wedge e_{k+1,n})$
- Successor state axioms : (if something became true, an action must have caused it)
  - $\neg l_k \wedge l_{k+1} \rightarrow (a_{k,1} \vee \dots \vee a_{k,j})$
- Exclusion axiom : exactly one action at a time
  - $a_{k,1} \vee \dots \vee a_{k,p}$  for each  $k$
  - $\neg a_{k,i} \vee \neg a_{k,j}$  for each  $k, i, j$

# SATPlan as CNF

$$O(C, F, 1) \wedge \neg I(B1, F, 1) \wedge \neg I(B2, F, 1),$$

$$O(C, F, 5) \wedge I(B1, F, 5) \wedge I(B2, F, 5),$$

$$\neg PC_k \vee (\neg O(C, F, k) \wedge O(C, F, k + 1))$$

$$\neg RC_k \vee (O(C, F, k) \wedge \neg O(C, F, k + 1))$$

$$\neg I1_k \vee (\neg O(C, F, k) \wedge \neg I(B1, F, k) \wedge I(B1, F, k + 1))$$

$$\neg I2_k \vee (\neg O(C, F, k) \wedge \neg I(B2, F, k) \wedge I(B2, F, k + 1))$$

$$(O(C, F, k) \vee \neg O(C, F, k + 1)) \vee (PC_k)$$

$$(\neg O(C, F, k) \vee O(C, F, k + 1)) \vee (RC_k)$$

$$(I(B1, F, k) \vee \neg I(B1, F, k + 1)) \vee (I1_k)$$

$$(\neg I(B1, F, k) \vee I(B1, F, k + 1))$$

$$(I(B2, F, k) \vee \neg I(B2, F, k + 1)) \vee (I2_k)$$

$$(\neg I(B2, F, k) \vee I(B2, F, k + 1)),$$

$$\neg RC_k \vee \neg PC_k$$

$$\neg RC_k \vee \neg O1_k$$

$$\neg RC_k \vee \neg O2_k$$

$$\neg PC_k \vee \neg O1_k$$

$$\neg PC_k \vee \neg O2_k$$

$$\neg O1_k \vee \neg O2_k,$$



# SATPlan

- Solutions

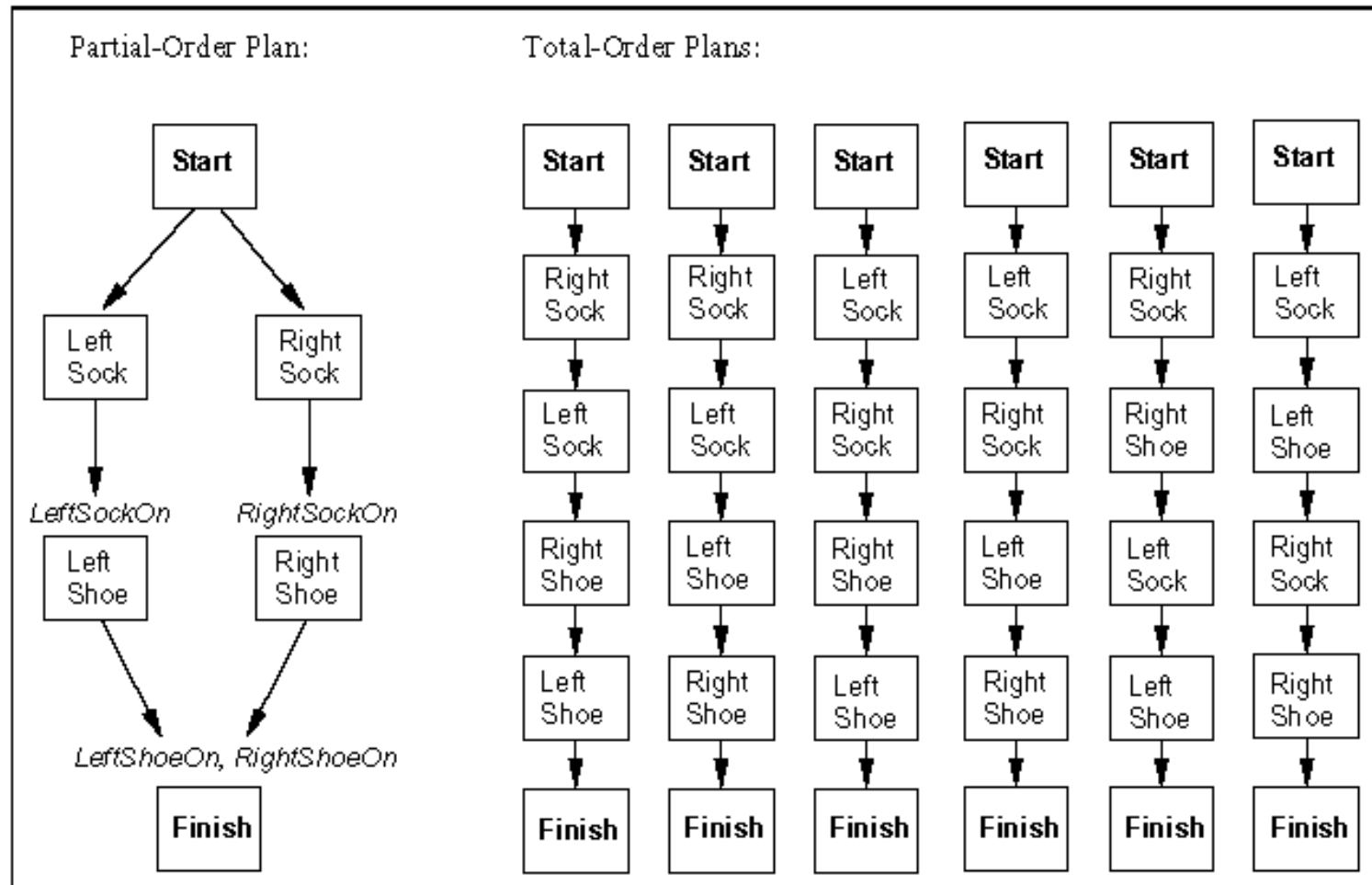
$RC_1$ ,  $IB1_2$ ,  $IB2_3$ , and  $PC_4$ .

$RC_1$ ,  $IB2_2$ ,  $IB1_3$ , and  $PC_4$ .

# Partial order planning

- Least commitment planning
- Nonlinear planning
- Search in the space of partial plans
- A state is a partial incomplete partially ordered plan
- Operators transform plans to other plans by:
  - Adding steps
  - Reordering
  - Grounding variables
- SNLP: Systematic Nonlinear Planning (McAllester and Rosenblitt 1991)
- NONLIN (Tate 1977)

# A partial order plan for putting shoes and socks



**Figure 11.6** A partial-order plan for putting on shoes and socks, and the six corresponding linearizations into total-order plans.

# Summary: Planning

- STRIPS Planning
- Situation Calculus
- Forward and backward planning
- Planning graph and GraphPlan
- SATplan
- Partial order planning
- Readings: RN chapter 10